Embedded Linux Conference Europe, October 2019

Offloading Network Traffic Classification

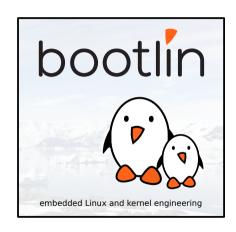
Maxime Chevallier

maxime.chevallier@bootlin.com

© Copyright 2004-2019, Bootlin.

Creative Commons BY-SA 3.0 license.

Corrections, suggestions, contributions and translations are welcome!





- Linux kernel engineer at Bootlin.
 - Linux kernel and driver development, system integration, boot time optimization, consulting. . .
 - Embedded Linux, Linux driver development, Yocto Project & OpenEmbedded and Buildroot training, with materials freely available under a Creative Commons license.
 - https://bootlin.com
- Contributions:
 - Worked on network (MAC, PHY, switch) engines.
 - Contributed to the Marvell EBU SoCs upstream support.
 - Also worked on SPI and real-time topics.



Preamble - goals

- ▶ Discover the classification operations in the kernel.
- Discover the hardware technologies used to offload packet classification
- Learn about the use cases for classification
- ▶ Based on **PPv2**'s behaviour and design, similar on other NICs



Introduction to Ingress Classification

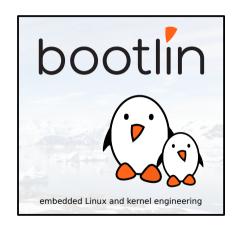
Maxime Chevallier

maxime.chevallier@bootlin.com

© Copyright 2004-2019, Bootlin.

Creative Commons BY-SA 3.0 license.

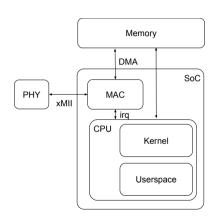
Corrections, suggestions, contributions and translations are welcome!





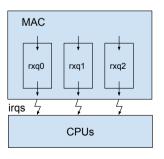
Packet path from the hardware to userspace

- 1. A frame arrives to the PHY
- 2. It is transferred to the MAC
- 3. The MAC performs offloaded operations
- 4. The packet is copied to RAM via DMA
- 5. The MAC raises an interrupt





- Upon receiving a frame, the packet goes through a Packet Processor
- ▶ The MAC receives the frames and places them into a buffer using DMA
- ▶ Descriptors for that buffer are placed into a **receive queue**
- ▶ Once the frame is received, the MAC raises an **interrupt**
- Receive queues can have dedicated interrupts, pinned to CPUs



- ▶ The interrupt is handled, mostly in softirg context
- napi is used to coalesce interrupts
- Packet processing is done on the CPU that handled the interrupt
- ▶ Before going up the network stack, the packet goes through the TC subsystem
- ▶ L2 handling, to deal with MAC filtering and VLANs
- L3 handling, to deal with routing
- ▶ L4 handling, where we find the **socket** that will consume the payload



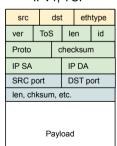
Classification

- Classification consists in identifying packets of interest
- We can then perform actions on these packets
- ▶ We first need to **dissect** the packet
- Determining the various attributes of interest isn't straightforward
- ► All fields don't have a fixed offset in the packet

VLAN, IPv4, TCP



IPv4, TCP



VLAN, IPv6, TCP



- ▶ A **flow** characterizes a group of packets that have a common source and destination.
- ▶ We group packets based on common attributes, such as :
 - ► The source and destination IP addresses (2-tuple)
 - ► The L4 protocol, source and destination ports (5-tuple)
- We manipulate flows to avoid reordering and optimize locality
- ▶ We need to extract the required information from the headers.

TC flower

- ► Traffic Control
- Used for traffic shaping, scheduling, policing and filtering
- In our case, we'll focus on the tc flower ingress filter
- **tc flower** is a **classifier**, which uses either software or hardware
- ► tc qdisc add dev eth0 ingress
- ▶ tc filter add dev ethO protocol ip parent ffff: flower ip_proto tcp dst_port 80 action drop

ethtool

- ethtool is used to interact with network drivers
- ▶ ethtool -N can be used to configure **n-tuple** filters
- It acts on specific flow types: tcp4, udp6, ether, etc.
- Rules are ordered, the first one that matches takes precedence
- ▶ ethtool -N eth0 flow-type tcp4 dst-port 80 action -1 loc 0
- Actions can be :
 - Steer to a Receive Queue
 - Steer to a RSS context
 - Drop



Embedded Linux Conference Europe, October 2019

Offloading Classification

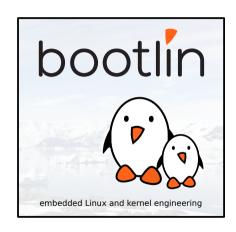
Maxime Chevallier

maxime.chevallier@bootlin.com

© Copyright 2004-2019, Bootlin.

Creative Commons BY-SA 3.0 license.

Corrections, suggestions, contributions and translations are welcome!



When and why

- ► Reduce CPU load
- ▶ Spread traffic across CPUs with per-cpu interrupts
- ► Early drop in case of Denial-of-Service attack
- Early redirection with switches

We must however be careful:

- ▶ The kernel might not see important packets
- The kernel might want to have access to the first packet of new flows
- Counters are not up to date anymore

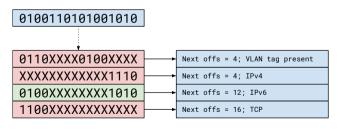
- ▶ Need to extract the required fields from the headers at wire speed
- ► These fields aren't always at a know position
- We need fast ways to lookup these fields, using a parser
- ▶ The attributes extracted by the parser are then used for **classification**





Ternary Content Addressable Memory

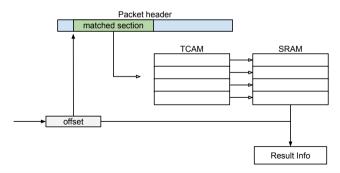
- Very fast lookups, but takes place on the die
- Addressed by value, returns the index of the first match
- Match on a ternary value : 0, 1 and X
- The matched pattern is extracted from the header starting from an offset
- The returned index is used to lookup a SRAM containing match actions





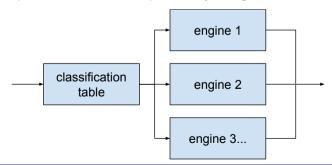
Parser

- Acts as a dissector
- Extract useful information from the packet header
- ▶ Take into account the various offsets due to DSA, VLAN and L3/L4
- Used as a pre-step for classification
- Often hardcoded in a firmware or a driver
- Multiple iterations per packet, flags are accumulated





- Uses information from the parser
- ▶ Can use several engines to classify and perform actions :
 - ► TCAM engines, for exact matches
 - Hash-based engines, for rate limiting and RSS
 - Logic engines for complex rules
- ▶ A final policing step decides what to do based on results from engines
- ▶ Not all these possibilities can be expressed by the generic frameworks





Receive Side Scaling

- Spread traffic across multiple CPUs
- Compute a hash from specific fields from the header
 - **s** : Source IP, **d** : Destination IP
 - **f** : Source port, **n** : Destination port
 - v : VLAN tag, m : Destination MAC
- Make sure that traffic from the same flows ends up on the same CPU
- Spreading is configured using an RSS Table
- ▶ ethtool -N eth2 rx-flow-hash tcp4 sdfn
- ▶ ethtool -X eth0 weight 2 1 1 0

PPv2 example

- PPv2 is found on Marvell SoCs, such as the Armada 70xx and 80xx
- Has a TCAM parser with 256 entries, performing up to 16 matches on 11B
- ▶ Classifier has one 512 instruction table, subdivided in subflows
- ► Has 4 classification engines :
 - ▶ C2 : TCAM match engine, 8B keys, 256 entries
 - ▶ C3 : Exact match engine, 12B keys, 4K entries
 - ▶ C4 : Classification and Marking engine, uses if-then-else constructs
 - **C3Hx**: Computes hashes, for RSS and C3 lookups.
- Can perform drop (in parser or classifier), steer to queue or RSS, limit traffic, modify and redirect packets.
- Parser and Classifier is shared between multiple ports



PPv2 : Current support

- Support for basic RSS
- Support steering on 2-tuple, 5-tuple and VLAN tag
- MAC and VLAN filtering, performed by the parser
- Support steering to RSS tables
- ▶ All Parser and Classifier configuration is done by the kernel, no firmware involved
- ▶ Only C2 and C3Hx engines are used, others are way too complex

- Offloading classification requires a lot of hardware configuration
- Most of the time, we need to limit ourselves to a subset of what the HW can do
- There are ongoing efforts to solve the issue of stats reporting
- ▶ Performance and power consumption improvements make it worth it
- ▶ In most cases, a firmware is in charge of configuring most of the tables

Thank you! Questions? Comments?

Maxime Chevallier — maxime.chevallier@bootlin.com

Slides under CC-BY-SA 3.0